CS 422/522  Design & Implementation of Operating Systems

Lecture 14: Cache & Virtual Memory

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Acknowledgement: some slides are taken from previous versions of the CS422/522 lectures taught by Prof. Bryan Ford and Dr. David Wolinsky, and also from the official set of slides accompanying the OSPP textbook by Anderson and Dahlin.

Definitions

♦ Cache
  - Copy of data that is faster to access than the original
  - Hit: if cache has copy
  - Miss: if cache does not have copy
♦ Cache block
  - Unit of cache storage (multiple memory locations)
♦ Temporal locality
  - Programs tend to reference the same memory locations multiple times
  - Example: instructions in a loop
♦ Spatial locality
  - Programs tend to reference nearby locations
  - Example: data in a loop
Cache concept (read)

Diagram showing the process of reading from cache:

1. Fetch Address
2. Check if Address is in Cache?
   - Yes: Store Value in Cache
   - No: Fetch Address

Cache concept (write)

Diagram showing the process of writing to cache:

1. Store Value at Address
2. Check if Address is in Cache?
   - Yes: If Write Through, Store Value in Cache
   - No: Fetch Address

Write through: changes sent immediately to next level of storage
Write back: changes stored in cache until cache block is replaced
**Memory hierarchy**

<table>
<thead>
<tr>
<th>Cache</th>
<th>Hit Cost</th>
<th>Size</th>
</tr>
</thead>
<tbody>
<tr>
<td>1st level cache/first level TLB</td>
<td>1 ns</td>
<td>64 KB</td>
</tr>
<tr>
<td>2nd level cache/second level TLB</td>
<td>4 ns</td>
<td>256 KB</td>
</tr>
<tr>
<td>3rd level cache</td>
<td>12 ns</td>
<td>2 MB</td>
</tr>
<tr>
<td>Memory (DRAM)</td>
<td>100 ns</td>
<td>10 GB</td>
</tr>
<tr>
<td>Data center memory (DRAM)</td>
<td>100 μs</td>
<td>100 TB</td>
</tr>
<tr>
<td>Local non-volatile memory</td>
<td>100 μs</td>
<td>100 GB</td>
</tr>
<tr>
<td>Local disk</td>
<td>10 ms</td>
<td>1 TB</td>
</tr>
<tr>
<td>Data center disk</td>
<td>10 ms</td>
<td>100 PB</td>
</tr>
<tr>
<td>Remote data center disk</td>
<td>200 ms</td>
<td>1 XB</td>
</tr>
</tbody>
</table>

i7 has 8MB as shared 3rd level cache; 2nd level cache is per-core

**Main points**

- Can we provide the illusion of near infinite memory in limited physical memory?
  - Demand-paged virtual memory
  - Memory-mapped files

- How do we choose which page to replace?
  - FIFO, MIN, LRU, LFU, Clock

- What types of workloads does caching work for, and how well?
  - Spatial/temporal locality vs. Zipf workloads
Hardware address translation is a power tool

- Kernel trap on read/write to selected addresses
  - Copy on write
  - Fill on reference
  - Zero on use
  - Demand paged virtual memory
  - Memory mapped files
  - Modified bit emulation
  - Use bit emulation

Demand paging (before)

<table>
<thead>
<tr>
<th>Page Table</th>
<th>Physical Memory</th>
<th>Disk</th>
</tr>
</thead>
<tbody>
<tr>
<td>Frame</td>
<td>Access</td>
<td></td>
</tr>
<tr>
<td>Virtual Page B</td>
<td>Frame for B Invalid</td>
<td>Page A</td>
</tr>
<tr>
<td>Virtual Page A</td>
<td>Frame for A R/W</td>
<td></td>
</tr>
</tbody>
</table>
**Demand paging (after)**

<table>
<thead>
<tr>
<th>Frame</th>
<th>Access</th>
</tr>
</thead>
<tbody>
<tr>
<td>Frame for B</td>
<td>R/W</td>
</tr>
<tr>
<td>Frame for A</td>
<td>Invalid</td>
</tr>
</tbody>
</table>

**Disk**

- Page A
- Page B

1. TLB miss
2. Page table walk
3. Page fault (page invalid in page table)
4. Trap to kernel
5. Convert virtual address to file + offset
6. Allocate page frame
   - Evict page if needed
7. Initiate disk block read into page frame
8. Disk interrupt when DMA complete
9. Mark page as valid
10. Resume process at faulting instruction
11. TLB miss
12. Page table walk to fetch translation
13. Execute instruction
Demand paging on MIPS (software TLB)

1. TLB miss
2. Trap to kernel
3. Page table walk
4. Find page is invalid
5. Convert virtual address to file + offset
6. Allocate page frame
   - Evict page if needed
7. Initiate disk block read into page frame
8. Disk interrupt when DMA complete
9. Mark page as valid
10. Load TLB entry
11. Resume process at faulting instruction
12. Execute instruction

Allocating a page frame

- Select old page to evict
- Find all page table entries that refer to old page
  - If page frame is shared
- Set each page table entry to invalid
- Remove any TLB entries
  - Copies of now invalid page table entry
- Write changes on page back to disk, if necessary
How do we know if page has been modified?

- Every page table entry has some bookkeeping
  - Has page been modified?
    * Set by hardware on store instruction
    * In both TLB and page table entry
  - Has page been recently used?
    * Set by hardware in page table entry on every TLB miss

- Bookkeeping bits can be reset by the OS kernel
  - When changes to page are flushed to disk
  - To track whether page is recently used

Keeping track of page modifications (before)

<table>
<thead>
<tr>
<th>TLB</th>
<th>Page Table</th>
<th>Physical Memory</th>
<th>Disk</th>
</tr>
</thead>
<tbody>
<tr>
<td>Frame</td>
<td>Access</td>
<td>Dirty</td>
<td>Page Frames</td>
</tr>
<tr>
<td>R/W</td>
<td>No</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Virtual Page B
- Frame for B
- Invalid

Virtual Page A
- Frame for A
- R/W
- No

Virtual Page A
- Frame A
- R/W
- No
Keeping track of page modifications (after)

TLB
<table>
<thead>
<tr>
<th>Frame</th>
<th>Access</th>
<th>Dirty</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>R/W</td>
<td>Yes</td>
</tr>
</tbody>
</table>

Page Table
<table>
<thead>
<tr>
<th>Frame</th>
<th>Access</th>
<th>Dirty</th>
</tr>
</thead>
<tbody>
<tr>
<td>Frame for B</td>
<td>Invalid</td>
<td></td>
</tr>
</tbody>
</table>

Physical Memory
<table>
<thead>
<tr>
<th>Page Frames</th>
</tr>
</thead>
</table>

Disk
| Old Page A |
| Old Page B |

Virtual Page B

Virtual Page A
<table>
<thead>
<tr>
<th>Frame</th>
<th>Access</th>
<th>Dirty</th>
</tr>
</thead>
<tbody>
<tr>
<td>Frame for A</td>
<td>R/W</td>
<td>Yes</td>
</tr>
</tbody>
</table>

Virtual or physical dirty/use bits

- Most machines keep dirty/use bits in the page table entry
- Physical page is
  - modified if any page table entry that points to it is modified
  - recently used if any page table entry that points to it is recently used
- On MIPS, simpler to keep dirty/use bits in the core map
  - Core map: map of physical page frames
Emulating a modified bit (Hardware Loaded TLB)

- Some processor architectures do not keep a modified bit per page
  - Extra bookkeeping and complexity
- Kernel can emulate a modified bit:
  - Set all clean pages as read-only
  - On first write to page, trap into kernel
  - Kernel sets modified bit, marks page as read-write
  - Resume execution
- Kernel needs to keep track of both
  - Current page table permission (e.g., read-only)
  - True page table permission (e.g., writeable, clean)

Emulating a recently used bit (Hardware Loaded TLB)

- Some processor architectures do not keep a recently used bit per page
  - Extra bookkeeping and complexity
- Kernel can emulate a recently used bit:
  - Set all recently unused pages as invalid
  - On first read/write, trap into kernel
  - Kernel sets recently used bit
  - Marks page as read or read/write
- Kernel needs to keep track of both
  - Current page table permission (e.g., invalid)
  - True page table permission (e.g., read-only, writeable)
Models for application file I/O

- Explicit read/write system calls
  - Data copied to user process using system call
  - Application operates on data
  - Data copied back to kernel using system call

- Memory-mapped files
  - Open file as a memory segment
  - Program uses load/store instructions on segment memory, implicitly operating on the file
  - Page fault if portion of file is not yet in memory
  - Kernel brings missing blocks into memory, restarts process

Advantages to memory-mapped files

- Programming simplicity, esp for large files
  - Operate directly on file, instead of copy in/copy out

- Zero-copy I/O
  - Data brought from disk directly into page frame

- Pipelining
  - Process can start working before all the pages are populated

- Interprocess communication
  - Shared memory segment vs. temporary file
From memory-mapped files to demand-paged virtual memory

- Every process segment backed by a file on disk
  - Code segment -> code portion of executable
  - Data, heap, stack segments -> temp files
  - Shared libraries -> code file and temp data file
  - Memory-mapped files -> memory-mapped files
  - When process ends, delete temp files

- Unified memory management across file buffer and process memory

Cache replacement policy

- On a cache miss, how do we choose which entry to replace?
  - Assuming the new entry is more likely to be used in the near future
  - In direct mapped caches, not an issue!

- Policy goal: reduce cache misses
  - Improve expected case performance
  - Also: reduce likelihood of very poor performance
A simple policy

- Random?
  - Replace a random entry

- FIFO?
  - Replace the entry that has been in the cache the longest time
  - What could go wrong?

FIFO in action

Worst case for FIFO is if program strides through memory that is larger than the cache
MIN, LRU, LFU

- **MIN**
  - Replace the cache entry that will not be used for the longest time into the future
  - Optimality proof based on exchange: if evict an entry used sooner, that will trigger an earlier cache miss

- **Least Recently Used (LRU)**
  - Replace the cache entry that has not been used for the longest time in the past
  - Approximation of MIN

- **Least Frequently Used (LFU)**
  - Replace the cache entry used the least often (in the recent past)

---

LRU/MIN for sequential scan

<table>
<thead>
<tr>
<th>LRU</th>
<th>MIN</th>
</tr>
</thead>
<tbody>
<tr>
<td>Reference</td>
<td>A</td>
</tr>
<tr>
<td>1</td>
<td>A</td>
</tr>
<tr>
<td>2</td>
<td>B</td>
</tr>
<tr>
<td>3</td>
<td>C</td>
</tr>
<tr>
<td>4</td>
<td>D</td>
</tr>
</tbody>
</table>
More page frames → fewer page faults?

- Consider the following reference string with 3 page frames
  - FIFO replacement
  - A, B, C, D, A, B, E, A, B, C, D, E
  - 9 page faults!

- Consider the same reference string with 4 page frames
  - FIFO replacement
  - A, B, C, D, A, B, E, A, B, C, D, E
  - 10 page faults

- This is called Belady’s anomaly
Belady's anomaly (cont'd)

Clock algorithm

- Approximate LRU
- Replace some old page, not the oldest unreferenced page
- Arrange physical pages in a circle with a clock hand
  - Hardware keeps “use bit” per physical page frame
  - Hardware sets “use bit” on each reference
  - If “use bit” isn’t set, means not referenced in a long time
- On page fault:
  - Advance clock hand
  - Check “use bit”
  - If “1” clear, go on
  - If “0”, replace page
Nth chance: Not Recently Used

- Instead of one bit per page, keep an integer
  - notInUseSince: number of sweeps since last use

- Periodically sweep through all page frames

```java
if (page is used) {
    notInUseSince = 0;
} else if (notInUseSince < N) {
    notInUseSince++;
} else {
    reclaim page;
}
```

Implementation note

- Clock and Nth Chance can run synchronously
  - In page fault handler, run algorithm to find next page to evict
  - Might require writing changes back to disk first

- Or asynchronously
  - Create a thread to maintain a pool of recently unused, clean pages
  - Find recently unused dirty pages, write mods back to disk
  - Find recently unused clean pages, mark as invalid and move to pool
  - On page fault, check if requested page is in pool!
  - If not, evict that page
Recap

- MIN is optimal
  - replace the page or cache entry that will be used farthest into the future
- LRU is an approximation of MIN
  - For programs that exhibit spatial and temporal locality
- Clock/Nth Chance is an approximation of LRU
  - Bin pages into sets of “not recently used”

How many pages allocated to each process?

- Each process needs minimum number of pages.

- Example: IBM 370 - 6 pages to handle SS MOVE instruction:
  - instruction is 6 bytes, might span 2 pages.
  - 2 pages to handle from.
  - 2 pages to handle to.

- Two major allocation schemes.
  - fixed allocation
  - priority allocation
Fixed allocation

- Equal allocation - e.g., if 100 frames and 5 processes, give each 20 pages.
- Proportional allocation - Allocate according to the size of process.

\[ s_i = \text{size of process } p_i \]
\[ S = \sum s_i \]
\[ m = \text{total number of frames} \]
\[ a_i = \text{allocation for } p_i = \frac{s_i}{S} \times m \]

\[ m = 64 \]
\[ s_1 = 10 \]
\[ s_2 = 127 \]
\[ a_1 = \frac{10}{137} \times 64 \approx 5 \]
\[ a_2 = \frac{127}{137} \times 64 \approx 59 \]

Priority allocation

- Use a proportional allocation scheme using priorities rather than size.
- If process \( P_i \) generates a page fault,
  - select for replacement one of its frames.
  - select for replacement a frame from a process with lower priority number.
Global vs. local allocation

- Global replacement – process selects a replacement frame from the set of all frames; one process can take a frame from another.

- Local replacement – each process selects from only its own set of allocated frames.

What to do when not enough memory?

- Thrashing: processes on system require more memory than it has.

* Each time one page is brought in, another page, whose contents will soon be referenced, is thrown out.
* Processes will spend all of their time blocked, waiting for pages to be fetched from disk
* I/O devices at 100% utilization but system not getting much useful work done

- What we wanted: virtual memory the size of disk with access time of physical memory
- What we have: memory with access time = disk access
Thrashing

- Process(es) “frequently” reference page not in mem
  - Spend more time waiting for I/O than getting work done
- Three different reasons
  - process doesn’t reuse memory, so caching doesn’t work
  - process does reuse memory, but it does not “fit”
  - individually, all processes fit and reuse memory, but too many for system.
- Which can we actually solve?

Making the best of a bad situation

- Single process thrashing?
  - If process does not fit or does not reuse memory, OS can do nothing except contain damage.

- System thrashing?
  - If thrashing arises because of the sum of several processes then adapt:
    * figure out how much memory each process needs
    * change scheduling priorities to run processes in groups whose memory needs can be satisfied (shedding load)
    * if new processes try to start, can refuse (admission control)
Working set model

- Working Set: set of memory locations that need to be cached for reasonable cache hit rate
- Size of working set = the important threshold
- The size may change even during execution of the same program.

Cache working set

![Graph showing hit rate vs. cache size (KB)]
Phase change behavior

Question

- What happens to system performance as we increase the number of processes?
  - If the sum of the working sets > physical memory?
Zipf distribution

- Caching behavior of many systems are not well characterized by the working set model

- An alternative is the Zipf distribution
  - Popularity $\sim 1/k^c$, for kth most popular item, $1 < c < 2$
  - "frequency inversely proportional to its rank in the frequency table (e.g., frequency word in English natural language)
    * Rank 1: "the" 7% (69,971 out of 1 million in "Brown Corpus")
    * Rank 2: "of" 3.5% (36,411 out of 1 million)
    * Rank 3: "and" 2.3% (28,852 out of 1 million)
Zipf examples

- Web pages
- Movies
- Library books
- Words in text
- Salaries
- City population
- ...

Common thread: popularity is self-reinforcing

Zipf and caching

Increasing the cache size continues to improve cache hit rates, but with diminishing returns
Cache lookup: fully associative

```
<table>
<thead>
<tr>
<th>address</th>
<th>value</th>
</tr>
</thead>
<tbody>
<tr>
<td>=?</td>
<td></td>
</tr>
<tr>
<td>=?</td>
<td></td>
</tr>
<tr>
<td>=?</td>
<td></td>
</tr>
<tr>
<td>=?</td>
<td></td>
</tr>
<tr>
<td>=?</td>
<td></td>
</tr>
</tbody>
</table>
```

match at any address?

```
yes
```

return value

Cache lookup: direct mapped

```
<table>
<thead>
<tr>
<th>address</th>
<th>value</th>
</tr>
</thead>
<tbody>
<tr>
<td>hash(address)</td>
<td></td>
</tr>
</tbody>
</table>
```

match at hash(address)?

```
yes
```

return value
Cache lookup: set associative

<table>
<thead>
<tr>
<th>hash(address)</th>
<th>address</th>
<th>value</th>
<th>match at hash(address)?</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>0x0053</td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

=? match at hash(address)?

=?

↓ yes

return value


Page coloring

- What happens when cache size >> page size?
  - Direct mapped or set associative
  - Multiple pages map to the same cache line

- OS page assignment matters!
  - Example: 8MB cache, 4KB pages
  - 1 of every 2K pages lands in same place in cache

- What should the OS do?
Page coloring

Processors → Virtual Address → Address Mod K → Cache → Memory

Memory: 0, K, 2K, 3K