



CS 428/528 Lecture 2

Logical Foundations & Coq

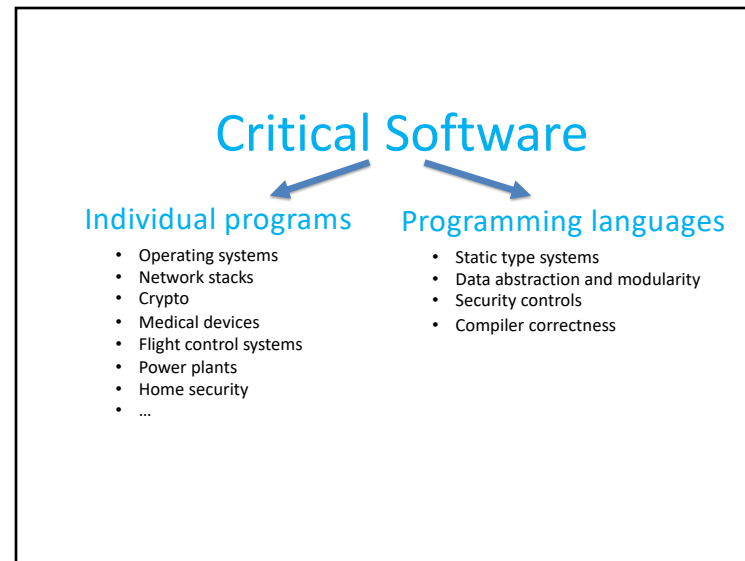
Zhong Shao
January 18, 2024

(Slides based on those from the Software Foundations
course material developed by Benjamin Pierce at Penn)

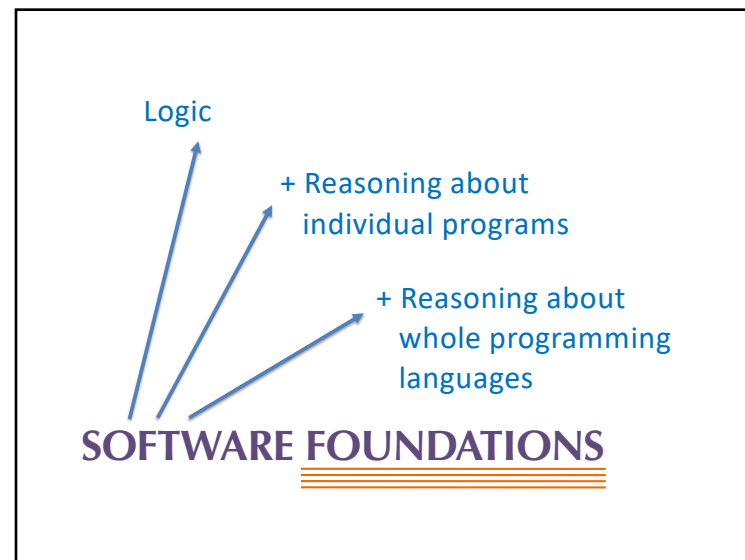
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How do we build software?
^
that works
^
(and be convinced
that it does)

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LOGICAL FOUNDATIONS

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Q: How do we know something is true?

A: We prove it

Q: How do we know that we have a *proof*?

A: We need to define what it means for something to be a proof.

A proof is a logical sequence of arguments, starting from some initial assumptions

Q: How do we know that we have a *valid* sequence of arguments? Can any sequence be a proof? E.g.

All humans are mortal

All Greeks are human

Therefore I am a Greek!

A: No, no, no! We need to think harder about valid ways of reasoning...



Aristotle
384 – 322 BC



Euclid
~300 BC

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First we need a language...

- **Gottlob Frege:** a German mathematician who started in geometry but became interested in logic and foundations of arithmetic.
- 1879 Published "*Begriffsschrift, eine der arithmetischen nachgebildete Formelsprache des reinen Denkens*" (Concept-Script: A Formal Language for Pure Thought Modeled on that of Arithmetic)
 - First rigorous treatment of functions and quantified variables
 - $\vdash A, \neg A, \forall x.F(x)$
 - First notation able to express arbitrarily complicated logical statements



Gottlob Frege
1848-1925



Images in this & following slides taken from Wikipedia.

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Formalization of Arithmetic

- 1884: *Die Grundlagen der Arithmetik* (The Foundations of Arithmetic)
- 1893: *Grundgesetze der Arithmetik* (Basic Laws of Arithmetic, Vol. 1)
- 1903: *Grundgesetze der Arithmetik* (Basic Laws of Arithmetic, Vol. 2)
- Frege's goals:
 - isolate logical principles of inference
 - derive laws of arithmetic from first principles
 - set mathematics on a solid foundation of logic

The plot thickens...

Just as Volume 2 was going to print in 1903,
Frege received a letter...

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Addendum to Frege's 1903 Book

*"Hardly anything more unfortunate can befall a scientific writer than to have one of the **foundations** of his edifice shaken after the work is finished. This was the position I was placed in by a letter of Mr. Bertrand Russell, just when the printing of this volume was nearing its completion."*

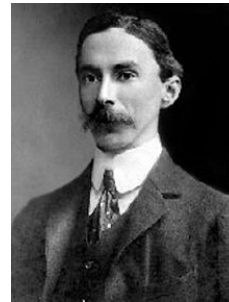
– Frege, 1903

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Bertrand Russell

- *Russell's paradox:*

1. Set comprehension notation:
 $\{x \mid P(x)\}$ "The set of x such that $P(x)$ "
2. Let X be the set (of sets) $\{Y \mid Y \notin Y\}$.
3. Ask the logical question:
Does $X \in X$ hold?
4. **Paradox!** If $X \in X$ then $X \notin X$.
If $X \notin X$ then $X \in X$.



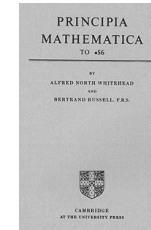
Bertrand Russell
1872 - 1970

- Frege's language could derive Russell's paradox \Rightarrow it was *inconsistent*.
- Frege's logical system could derive anything. Oops (!!)

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Aftermath of Frege and Russell

- Frege came up with a fix... but it made his logic trivial :-(
- 1908: Russell fixed the inconsistency of Frege's logic by developing a *theory of types*.
- 1910, 1912, 1913, (revised 1927): *Principia Mathematica* (Whitehead & Russell)
 - Goal: axioms and rules from which *all* mathematical truths could be derived.
 - It was a bit unwieldy...



"From this proposition it will follow, when arithmetical addition has been defined, that $1+1=2$."
 —Volume I, 1st edition, page 379

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Logic in the 1930s and 1940s

- 1931: Kurt Gödel's first and second incompleteness theorems.
 - Demonstrated that any consistent formal theory capable of expressing arithmetic cannot be complete.
 -
- 1936: Gödel proves *consistency* of arithmetic.
- 1936: Church introduces the λ -calculus.
- 1936: Turing introduces Turing machines
 - Is there a decision procedure for arithmetic?
 - Answer: no, it's undecidable
 - The famous "halting problem"
 - N.b.: Only in 1938 did Turing get his Ph.D.
- 1940: Church introduces the *simple theory of types*



Kurt Gödel
1906 - 1978



Gerhard Gentzen
1909 - 1945



Alonzo Church
1903 - 1995




Alan Turing
1912 - 1954

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Fast Forward...

- Two logicians in 1958 (Haskell Curry) and 1969 (William Howard) observe a remarkable correspondence:


types	~	propositions
programs	~	proofs
computation	~	simplification



Haskell Curry
1900 – 1982



William Howard
1926 –



N.G. de Bruijn
1918 - 2012


- 1967 – 1980's: N.G. de Bruijn runs Automath project
 - uses the Curry-Howard correspondence for computer-verified mathematics
- 1971: Jean-Yves Girard introduces System F
- 1972: Girard introduces F_ω
- 1972: Per Martin-Löf introduces intuitionistic type theory
- 1974: John Reynolds independently discovers System F

Basis for modern type systems:
OCaml, Haskell, Scala, Java, C#, ...


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... to the Present

- 1984: Coquand and Huet first begin implementing a new theorem prover "Coq"
- 1985: Coquand introduces the calculus of constructions
 - combines features from intuitionistic type theory and F_ω
- 1989: Coquand and Paulin extend CoC to the calculus of inductive constructions
 - adds "inductive types" as a primitive
- 1992: Coq ported to Xavier Leroy's OCaml
- 1990's: up to Coq version 6.2
- 2000-2015: up to Coq version 8.4
- 2017: Coq version 8.6



Thierry Coquand
1961 –



Gérard Huet
1947 –

} Too many contributors to list here...

- 2013: Coq receives ACM Software System Award

<http://coq.inria.fr/refman/Reference-Manual002.html>

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So much for foundations... what about the "software" part?

(LANGUAGE) PROGRAMMING FOUNDATIONS

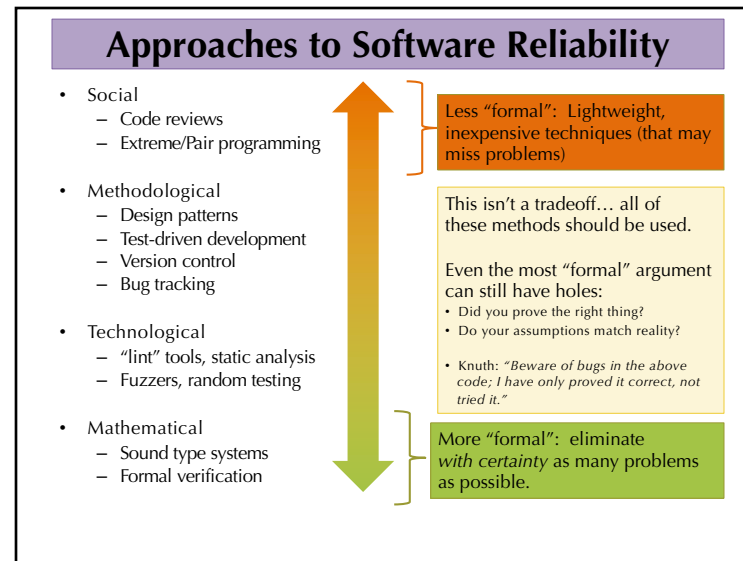
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Building Reliable Software

- Suppose you work at (or run) a software company.
- Suppose, like Frege, you've sunk 30+ person-years into developing the "next big thing":
 - Boeing Dreamliner2 flight controller
 - Autonomous vehicle control software for Nissan
 - Gene therapy DNA tailoring algorithms
 - Super-efficient green-energy power grid controller
- Suppose, like Frege, your company has invested a lot of material resources that are also at stake.
- How do you avoid getting a letter like the one from Russell?

Or, worse yet, *not* getting the letter,
with disastrous consequences down the road?

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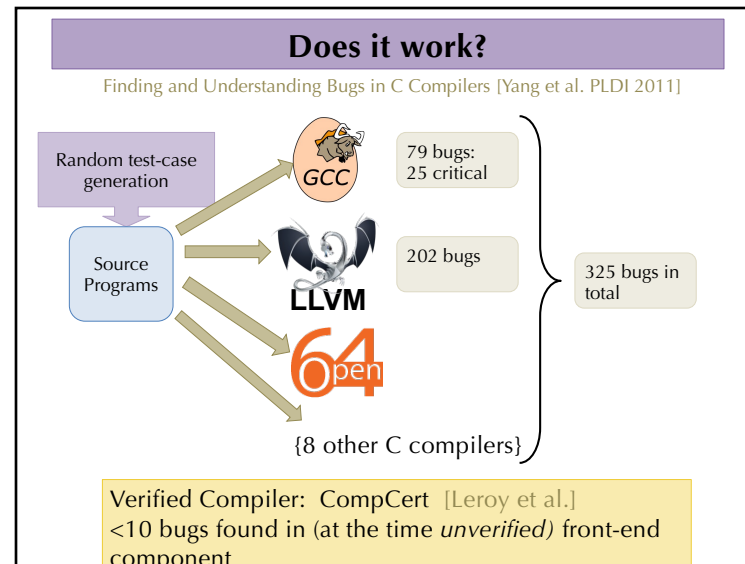
Can formal methods scale?

Use of formal methods to verify full-scale software systems is a hot research topic!

- **CompCert** – fully verified C compiler
Leroy, INRIA
- **Vellvm** – formalized LLVM IR
Zdancewic, Penn
- **Verified Software Toolchain**
Appel, Princeton
- **Bedrock** – web programming, packet filters
Chlipala, MIT
- **CertiKOS** – certified OS kernel
Shao, Yale

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Regehr's Group Concludes

The striking thing about our CompCert results is that the *middle-end bugs* we found in all other compilers are *absent*. As of early 2011, the under-development version of *CompCert is the only compiler we have tested for which Csmith cannot find wrong-code errors*. This is not for lack of trying: we have devoted about six CPU-years to the task. *The apparent unbreakability of CompCert supports a strong argument that developing compiler optimizations within a proof framework, where safety checks are explicit and machine-checked, has tangible benefits for compiler users.*

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